

# TREE-ORIENTED DATA ACQUISITION ARCHITECTURE

Vit NOVOTNY<sup>1</sup>

<sup>1</sup>Department of Telecommunication, Faculty of Electrical Engineering and Communication, Brno University of Technology, Purkynova 118, 612 00 Brno, Czech Republic

novotnyv@feec.vutbr.cz

**Abstract.** Hierarchical system of network nodes is suitable solution how to collect and how to pre-process data from large amount of end-nodes. By contrast to flat (one layer) architecture there are special intermediary nodes used and they are called summarization nodes. These special nodes have to be suitably placed in the network to enable efficient data collection and their number in the hierarchy is one of the key parameters of the architecture. The article deals with the tree architecture design, with its optimisation and with the problem of limited number of summarization nodes.

## Keywords

*Data acquisition, tree architecture, multicast, sensor networks, SSM, multimedia, RTP/ RTCP.*

## 1. Introduction

There are several ways of data collection and processing in the network environment. First model is called centralized, where there is only one data centre where all pieces of information are collected, processed and available. The second one is called hierarchical system created by a tree of servers with local data and using links among them the requested information can be found through the tree hierarchy. The third model is distributed model where the pieces of information are distributed among equivalent data centres and using one sophisticated directory services with one account requested information can be obtained.

Many applications are using plain centralized model. There is no problem with data acquisition and with data processing provided the number of data sources is fairly low and data flows are weak and low frequent. When these conditions are not fulfilled either the centre itself or data links to the centre can be overloaded or allowed data transmission frequency is very low. When

the data acquisition is auxiliary procedure of the service, the available bandwidth for such procedure is strictly limited and the situation becomes even worse. This is the case of applications like IP-TV where the main procedure of the service is the multimedia streaming using RTP protocol and the multicast transmission and the session quality parameter collection using RTCP protocol is an optional though useful supplementary service [1], [2], [3]. The transmission capacity of RTCP is limited for 5 % of total service bandwidth and it causes large delays in sending RTCP (feedback) data from each receiver for large-scale media streaming services based on Source-Specific Multicast (SSM), [7]. Similar problem arises also with other applications focused on data acquisition in the case of large-scale systems.

## 2. Hierarchical Data Acquisition System

To combat the problem the hierarchical system for data acquisition has been proposed in [5], [6] and modified in [7]. In addition to the data centre and data sources such tree contains special nodes called summarization nodes, see Fig. 1.

The data is periodically sent from data sources (terminals or sensors) to assigned summarization node. The summarization node aggregates data from a group of terminals of the size  $n_B$  and again periodically sends to assigned summarization node at the higher level. The summarization nodes are also organized into groups of size  $n_S$ . Structure of Receiver Summary Information (RSI) message was specified in [5]. The message includes sub-report blocks (SRB) that contain distribution information about particular features like a packet loss or a jitter.

To enable efficient transmission of information about the session from the data centre to the terminals an extension of original RTCP specification in the form of Extend Report (XR) message had to be adopted [3]. The

XR RTCP summarization packet consists basic information for the terminals mainly how to calculate the message transmission period. In the case of SSM (Source-Specific Multicast) service the message is sent in multicast manner so that together with the summarization method it decreases the overhead and saves the bandwidth.

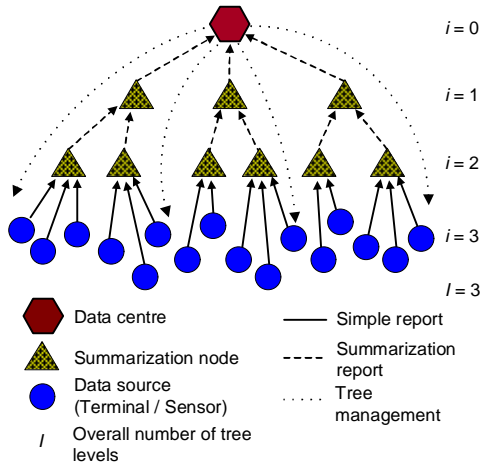


Fig. 1: Tree design for large-scale data acquisition.

Optionally in addition to the summarization process (when detailed information is lost) the summarization nodes can store detailed information obtained from terminals or lower level summarization nodes for some time period to allow the data centre to get detailed information about particular terminal or group of terminals when necessary.

Because of large group of terminals division into a big number of smaller groups the bandwidth restriction is not the problem and the message transmission period of the terminals remains fairly low even if the overall number of terminals rises. Especially this is the case of multimedia multicast sessions which can vary substantially in size. The overall delay that is bounded by the time instant when data is generated (or measured) in the terminal (sensor) and by the time instant when the data is received in the data centre consists of particular transmission delays between transmission instants of adjacent layers in the tree. The situation is depicted in Fig. 2:

When the tree consists of  $I$  layers, i.e.  $(I-1)$  summarization layers and one terminal layer, a formula for the overall delay  $T_R$  between data generation (measurement) and its reception in the data processing centre can be derived:

$$T_R = \tau_{MT} + \sum_{i=1}^{I-1} \tau_i, \quad (1)$$

provided the transport delay through the network is neglected. Variable  $\tau_{MT}$  is the delay between measurement (data generation) and transmission instants

and  $\tau_i$  is the delay between summarized message transmission instants at linked summarization nodes in adjacent layers.

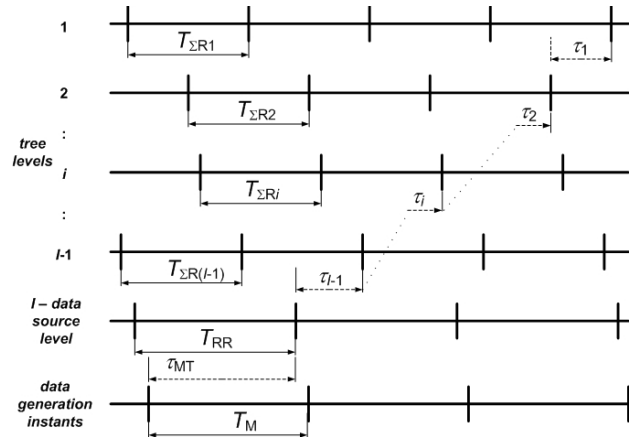


Fig. 2: Time instants of data generation and message transmissions in hierarchical data acquisition system.

The worst case for the delay will be when all summarization nodes at all levels of the tree and also the terminals (sensors) are synchronized, i.e. all of them transmit messages almost at the same time instants. Then the formula (1) will convert to

$$T_{RTW} = T_{RR} + \sum_{i=1}^{I-1} T_{\Sigma Ri}. \quad (2)$$

Provided the transmission periods are the same through whole tree the formula (2) changes into the form

$$T_{RTW} = T_{RR} + (I-1)T_{\Sigma R}, \quad (3)$$

where  $T_{RR}$  is the transmission period of the group of terminals (it depends on the number  $n_B$  of terminals in the group, message length and the allocated bandwidth, [7]),  $I$  is the number of levels in the tree (it depends on the total number of terminals, on the number of terminals in the group and on the number of summarization nodes in the group) and  $T_{\Sigma R}$  is the message transmission period of the summarization node group (it depends on the number  $n_S$  of summarization nodes in the group, summarization message length and the allocated bandwidth, [7]).

Now several problems come out. First group of problems are how to manage the tree when the number of terminals rises or declines, how to keep it in balanced form and how to minimize the total delay specified by (3). In addition to this the problem of the number of required summarization nodes should be addressed. At the beginning of our research we considered that the summarization nodes are only terminals with special functionality [7]. It was found that there would be lot of overhead with the management of such tree especially when the tree is variable in a large extent, i.e. the terminals will enter and leave the session frequently; this is the case of multimedia streaming sessions. Also this functionality would require additional power and energy

that is unwanted issue especially in the case of wireless terminals (sensors) with very limited computational power and energy. Therefore in later research ([8], [10]) the summarization nodes are considered as special nodes (or software modules) that are managed by the service provider. Such summarization nodes have higher computational power, larger storage capacity for temporary data and fixed location. This last feature is very important when tree structure is established according to the location of terminals, [10].

### 3. Tree Optimization

When the service provider intends to implement a service based on the tree architecture described above before implementation some initial conditions have to be considered: bandwidth (or maximum data flow) allocated for the data acquisition  $BW_A$  (it will be allocated for each group of terminals or summarization nodes), expected number of data sources (terminals)  $n_T$ , maximum period (or delay) of data collection  $T_{Rmax}$ , length  $PL_{RR}$  of plain messages generated by the terminals and the length  $PL_{\Sigma R}$  of summarization packets generated by summarization nodes. Additional constraints can be: maximum overall number of available summarization nodes  $N_{STmax}$ , minimum periods of message transmission in a group of terminals  $T_{RRmin}$  and in a group of summarization nodes  $T_{\Sigma Rmin}$  and some others. The goal is to find such tree that meets all of these conditions and restrictions.

Equation (3) shows how to calculate the largest overall delay  $T_R$  (and also the maximum time period of data acquisition) between data generation (measurement) in terminals (data sources) and its reception in the data processing centre. It can be worked out in more detailed form:

$$T_{RTW} = T_{RR} + (I-1)T_{\Sigma R} = \tau_R n_B + (I-1)\tau_{\Sigma} n_S = (BW_A)^{-1} [PL_{RR} n_B + (I-1)PL_{\Sigma R} n_S], \quad (4)$$

where  $I$  is a number of tree levels,  $\tau_R$  is a time interval consumed by one message send by a terminal and  $\tau_{\Sigma}$  is a time interval consumed by one summarization message generated by summarization node,  $n_B$  is the number of terminals in one group of terminals,  $n_S$  is the number of nodes in the group of summarization nodes (the rest of symbols are explained in the text above).

The number of levels with summarization nodes in the tree, i.e. the value  $(I-1)$ , can be calculated from the condition

$$n_S^{(I-2)} < \frac{n_T}{n_B} \leq n_S^{(I-1)}. \quad (5)$$

Then

$$\log_{n_S} \left( \frac{n_T}{n_B} \right) \leq (I-1) < \log_{n_S} \left( \frac{n_T}{n_B} \right) + 1. \quad (6)$$

As  $I$  is an integer number the nearest higher integer will be

$$(I-1) = \log_{n_S} \left( \frac{n_T}{n_B} \eta_I \right); \quad \eta_I \in \langle 1, n_S \rangle. \quad (7)$$

Then (4) changes into form

$$T_{RTW} = \tau_R n_B + (I-1)\tau_{\Sigma} n_S = \tau_R n_B + (\tau_{\Sigma} n_S) \left[ \log_{n_S} \left( \frac{n_T}{n_B} \eta_I \right) \right]. \quad (8)$$

so that the optimization process of delay minimization consists of seeking the numbers  $n_B$  and  $n_S$ . An example of the delay course according to (8) is shown in Fig. 3:

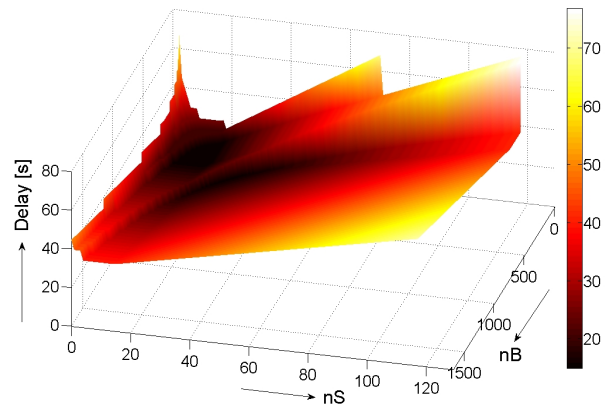


Fig. 3: Course of the worst-case total delay according to (8).

When discontinuous function (8) is replaced by continuous one (without correction parameter  $\eta_I$ ), we obtain expression

$$T_{RTW} = \tau_R n_B + (\tau_{\Sigma} n_S) \left[ \log_{n_S} \left( \frac{n_T}{n_B} \right) \right]. \quad (9)$$

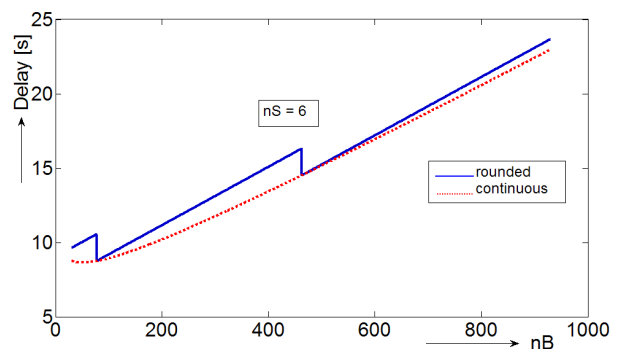


Fig. 4: Comparison of discontinuous (8) and continuous (9) representations of the worst-case delay.

The worst-case total delay values obtained from the optimization process with continuous function are quite close to and always a better than when

discontinuous function is considered (due to the fact that  $\eta_i \geq 1$ ), see Fig. 4:

The goal of optimization is to find its global extreme (minimum) in this region. Global extreme can be located either in local extremes of the function or at the boundary of definition domain. The function is continuous in whole region and smooth, therefore first and also second derivatives can be calculated and stationary points of the function can be found:

$$\frac{\partial T_{RTW}}{\partial n_B} = \tau_R - \tau_\Sigma \frac{n_S}{n_B \ln n_S}, \quad (10)$$

$$\frac{\partial T_{RTW}}{\partial n_S} = \tau_\Sigma \frac{\ln\left(\frac{n_T}{n_B}\right)}{\ln^2 n_S} (\ln n_S - 1). \quad (11)$$

Stationary points are the candidates for local extremes and they can be calculated from the conditions that first derivatives (10) and (11) are put equal zero and the results are:

$$n_{Ss1} = e \text{ (i.e. } 2.71828\dots) + \eta_S = 3 \quad (12)$$

and when non-rounded  $n_{Ss1}$  is used for  $n_{Bs1}$  calculations

$$n_{Bs1} = \frac{\tau_\Sigma}{\tau_R} \frac{n_{Ss1}}{\ln n_{Ss1}} = \frac{\tau_\Sigma}{\tau_R} e + \eta_B = \frac{PL_{\Sigma R}}{PL_{RR}} e + \eta_B, \quad (13)$$

$$\eta_B \in (-0.5; +0.5).$$

To prove, whether the local minimum was found, it is necessary to check sufficient conditions for the existence of local minimum

$$D_1 = \frac{\partial^2 T_{RTW}}{\partial n_B^2} \bigg|_{\substack{n_B=n_{Bs} \\ n_S=n_{Ss}}} \geq 0,$$

$$D_2 = \begin{vmatrix} \frac{\partial^2 T_{RTW}}{\partial n_B^2} & \frac{\partial^2 T_{RTW}}{\partial n_B \partial n_S} \\ \frac{\partial^2 T_{RTW}}{\partial n_B \partial n_S} & \frac{\partial^2 T_{RTW}}{\partial n_S^2} \end{vmatrix} \bigg|_{\substack{n_B=n_{Bs} \\ n_S=n_{Ss}}} > 0. \quad (14)$$

When the results (12) and (13) are used in (14) we get:

$$D_1 = \frac{\tau_R^2}{\tau_\Sigma e},$$

$$D_2 = \frac{\tau_R^2}{e^2} \left[ \ln\left(n_T \frac{\tau_R}{\tau_\Sigma}\right) - 1 \right]. \quad (15)$$

The inequality  $D_1 > 0$  is always met and the inequality  $D_2$  will be fulfilled when

$$n_T \frac{\tau_R}{\tau_\Sigma} > e. \quad (16)$$

Again in the example of RTCP presented in [7] the length of receiver report  $PL_{RR}$  was 736 bits and the length of summarization report  $PL_{\Sigma R}$  was 11296 bits. In the case when the same link bandwidths are assigned both to terminals and summarization nodes (16) has the form

$$n_T \frac{\tau_R}{\tau_\Sigma} = n_T \frac{PL_{RR}}{PL_{\Sigma R}} = n_T \frac{736}{11296} \approx 0.065 n_T > e, \quad (17)$$

$$n_T > 42.$$

This condition is quite easy to meet.

Provided the condition (16) is met the (12) and (13) specify local delay minimum:

$$T_{RTW1} \left( n_{Bs1} = \frac{\tau_\Sigma}{\tau_R} e + \eta_B, n_{Ss1} = 3 \right) =$$

$$= \tau_\Sigma \left[ e + 3 \log_3 \left( \frac{n_T}{\frac{\tau_\Sigma}{\tau_R} e + \eta_B} \right) \right] + \tau_R \eta_B. \quad (18)$$

When minimum periods of message transmission in a group of terminals  $T_{RRmin}$  and in a group of summarization nodes  $T_{\Sigma Rmin}$  are required then minimum values  $n_{Bmin}$ ,  $n_{Smin}$  are set:

$$n_{Bmin} = \frac{BW_A}{PL_{RR}} T_{RRmin} + \eta_{Bmin} = \frac{T_{RRmin}}{\tau_R} + \eta_{Bmin},$$

$$\eta_{Bmin} \in (0, 1), \quad (19)$$

$$n_{Smin} = \frac{BW_A}{PL_{\Sigma R}} T_{\Sigma Rmin} + \eta_{Smin} = \frac{T_{\Sigma Rmin}}{\tau_\Sigma} + \eta_{Smin},$$

$$\eta_{Smin} \in (0, 1).$$

Then the absolute minimum will be reached for the smallest  $n_S = n_{Smin}$  and for  $n_B = n_{Bmin}$ .

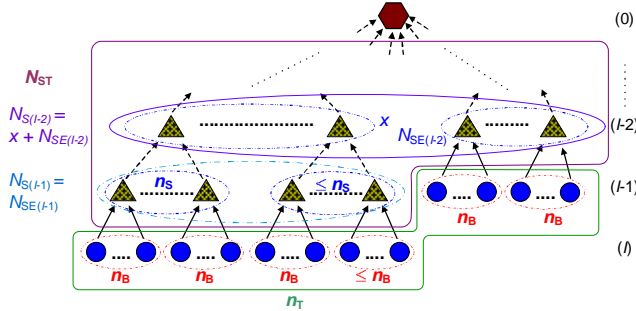
## 4. Summarization Nodes

The service provider has no unlimited number of summarization nodes available and therefore the overall number of required summarization nodes  $N_{ST}$  in the tree hierarchy is also very important parameter and should be optimized. The total number of summarization nodes can be calculated as follows (see Fig. 5):

$$N_{ST} = \sum_{i=1}^{I-1} N_{Si} = \sum_{i=1}^{I-2} n_S^i + N_{S(I-1)} = \frac{n_S^{I-2} - 1}{1 - n_S^{-1}} + N_{S(I-1)}, \quad (20)$$

where  $I$  is the number of levels in the tree,  $N_{Si}$  is the number of summarization nodes at the level  $i$ ,  $n_S$  is the number of summarization nodes in one group and  $N_{S(I-1)}$  is the number of summarization nodes at the level  $I-1$ .

The parameter  $n_s$  is known and therefore the task is to calculate the variable  $N_{S(I-1)}$ . As shown in Fig. 5: the terminals (sensors) are connected to the summarization nodes at two layers,  $(I-2)$  and  $(I-1)$  respectively.



**Fig. 5:** Tree architecture for data acquisition with highlighting of last three levels.

These summarization nodes can be called summarization endpoint nodes  $N_{SE}$  and its total number  $N_{SET}$  ( $T$  = total) can be expressed by formula

$$N_{SET} = N_{SE(I-2)} + N_{SE(I-1)}, \quad (21)$$

where  $N_{SE(I-2)}$  and  $N_{SE(I-1)}$  are endpoint summarization nodes at levels  $(I-2)$  and  $(I-1)$  respectively. As the  $(I-1)$  layer is the last layer of the summarization nodes it is clear that  $N_{S(I-1)} = N_{SE(I-1)}$ . The parameter  $N_{SET}$  is an integer figure and it can be calculated by the equation

$$N_{SET} = \frac{n_T}{n_{Bmax}} + \eta_E; \quad \eta_E \in (0, 1), \quad (22)$$

where  $n_{Bmax}$  is the maximum number of terminals in one group.

To obtain the total number of required summarization nodes  $N_{ST}$  it is necessary to calculate  $N_{SE(I-1)}$ . To get this parameter we need  $N_{SE(I-2)}$  first. When new terminals are to be added and the current tree is not sufficient, additional layer of summarization nodes has to be added. An appropriate number of summarization nodes  $x$  that will loose the terminals for next-layer summarization nodes (where the maximum of summarization nodes therefore can be  $x \cdot n_s$ ) can be calculated from (23):

$$\begin{aligned} N_{SET} &= N_{SE(I-2)} + N_{SE(I-1)} = \\ &= N_{S(I-2)} - x + N_{S(I-1)} = \\ &= n_s^{I-2} - x + N_{S(I-1)} \leq n_s^{I-2} - x + x n_s. \end{aligned} \quad (23)$$

The closest larger integer is

$$x = \frac{\frac{n_T}{n_{Bmax}} + \eta_E - n_s^{I-2}}{n_s - 1} + \eta_x; \quad \eta_x \in (0, 1). \quad (24)$$

Then the total number of required summarization nodes is:

$$\begin{aligned} N_{ST} &= \sum_{i=1}^{I-1} N_{Si} = \sum_{i=1}^{I-2} n_s^i + N_{S(I-1)} = \frac{n_s^{I-2} - 1}{1 - n_s^{-1}} + \\ &+ \frac{N_{SET} - n_s^{I-2}}{1 - n_s^{-1}} + \eta_x = \frac{N_{SET} - 1}{1 - n_s^{-1}} + \eta_x = \\ &= \frac{1}{1 - n_s^{-1}} \left[ \frac{n_T}{n_{Bmax}} - 1 \right] + \frac{\eta_E}{1 - n_s^{-1}} + \eta_x. \end{aligned} \quad (25)$$

The formula (25) shows that the overall number of summarization nodes is mainly influenced by total number of receivers  $n_T$  and by the number of terminals in one group  $n_{Bmax}$ . Parameter  $n_s$  does not have big impact on the  $N_{ST}$  when  $n_s \gg 1$ .

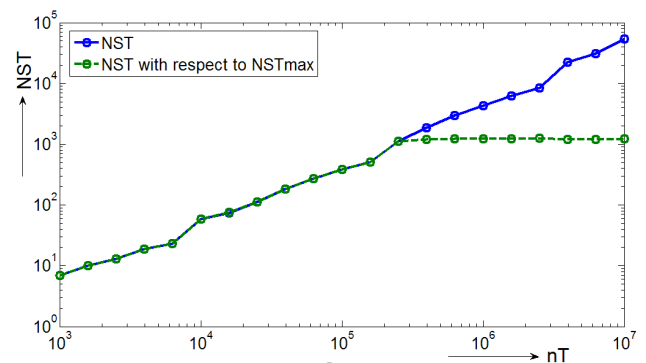
The rest of the formula (25), i.e. expression  $\frac{\eta_E}{1 - n_s^{-1}} + \eta_x$  will be always  $< 3$  (even when  $n_s = 2$ , which is the smallest value, and both  $\eta_x \rightarrow 1$  and  $\eta_E \rightarrow 1$ ). Then it holds that

$$N_{ST} \leq \frac{1}{1 - n_s^{-1}} \left[ \frac{n_T}{n_{Bmax}} - 1 \right] + 3. \quad (26)$$

If a provider has  $N_{STmax}$  summarization nodes available we can derive the minimum number of terminals in one group  $n_{Bmxm}$  from (26) and we get

$$n_{Bmxm} \geq \frac{n_T}{(N_{STmax} - 3)(1 - n_s^{-1}) + 1}. \quad (27)$$

When required parameter  $n_{Bmxm}$  obtained from (27) is smaller than optimum parameter  $n_{Bopt}$ , the point  $(n_{Bopt}, n_{Sopt})$  will be used as the best value for the number of terminals in one group and for summarization nodes in one group respectively. Otherwise the new optimum value  $n_{Bopt}$  larger than  $n_{Bmxm}$  and  $n_{Sopt}$  will be searched. The results are separately compared in Fig. 6: and in Fig. 7:



**Fig. 6:** The courses of required number of summarization nodes for different numbers of terminals ( $n_T$ ) without a and with respect to the demand on  $N_{STmax} = 1190$ .



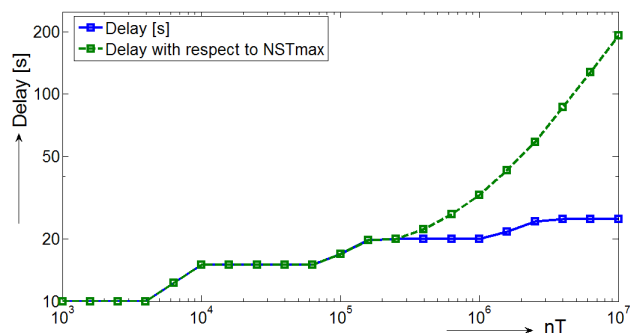


Fig. 7: The courses of minimum delay for different numbers of terminals (nT) without and with respect to the demand on NSTmax = 1190.

## 5. Conclusion

This article dealt with the problem of hierarchical data acquisition. The process of tree design was presented and some problems related to it were addressed like minimization of the total acquisition delay and the limited figure of summarization nodes. The delay optimum was found and tree parameters were derived. Influence of limited number of summarization nodes was considered and proved by simulations in Matlab environment. Separate paper will address the problem of end nodes (terminals) organisation according to their localities.

## Acknowledgements

This paper a the research results were got with the support of Ministry of Education, Youth and Sports the project No. MSM0021630513.

## References

- [1] SCHULZRINNE, H.; CASNER, S.; FREDERICK, R.; JACOBSON, V. *RTP: A Transport Protocol for Real-Time Applications*. Internet Draft, IETF RFC3550, 2003. pp. 1-73
- [2] HOLBROOK H.; CAIN, B. *Source-Specific Multicast for IP*. Internet Draft, IETF, 2004. pp. 1-17
- [3] BHATTACHARYYA, S. *An Overview of Source-Specific Multicast (SSM)*. Request for Comments 3569, IETF, 2003. pp. 1-14
- [4] CHESTERFIELD, J.; SCHOOLER, E., M. An extensible RTCP control framework for large multimedia distributions, In *Second IEEE International Symposium on Network Computing and Applications, 2003. NCA 2003*. 2003. pp. 351 – 358. ISBN 0-7695-1938-5.
- [5] CHESTERFIELD, J.; OTT, J.; SCHOOLER, E., M. *RTCP Extensions for Single-Source Multicast Sessions with Unicast Feedback*. Internet draft, IETF, 2007. pp. 1-40
- [6] EL-MARAKBY, R.; HUTCHISON, D. Scalability Improvement of the Real-Time Control Protocol (RTCP) Leading to Management Facilities in the Internet. In *Proceedings of Third IEEE Symposium on Computers & Communications, ISCC*. 1998. pp. 125 – 129. ISBN 0-8186-8538-7.
- [7] CASTRO, M.; DRUSCHEL, P.; KERMARREC, A., M.; ROWSTRON, A., I., T. A large-scale and decentralized application-level multicast infrastructure. *IEEE Journal on Selected Areas in Communications*. 2002. Vol. 20, No. 8, pp. 1489-1499. ISSN 0733-8716.
- [8] NOVOTNÝ, V.; KOMOSNÝ, D. Optimization of Large-Scale RTCP Feedback Reporting in Fixed and Mobile Networks. In *Proceedings of Third International Conference on Wireless and Mobile Communications, 2007. ICWMC '07*. March 2007, pp. 1-6. ISBN: 0-7695-2796-5.
- [9] NOVOTNÝ, V.; KOMOSNÝ, D. Large-Scale RTCP Feedback Optimization. *Journal of Networks*, 2008, Vol. 3, No. 3, pp. 1-10. ISSN 1796-2056.
- [10] KOMOSNÝ, D.; BURGET, R.; NOVOTNÝ, V. Tree Transmission Protocol for Feedback Distribution in IPTV Systems. In *Proceedings of the Seventh IASTED International Conference on Communication Systems and Networks*. Palma de Mallorca, Spain: International Association of Science and Technology for Development, 2008. pp. 1-7. ISBN 978-0-88986-758-1.

## About Author

**Vit NOVOTNY** was born in 1969. He received his M.Sc. from Brno University of Technology in 1992, his Ph.D. degree received at the Brno University of Technology in 2001 in Electronics and communication technologies and he became the assistant professor at the same university in 2005, also in the area of “electronics and communication technologies”. In the past he did the research in the areas of non-filtering applications of switched capacitors and of the current and voltage conveyors. Current professional interests are mobile and packet data networks, their services and terminal equipment. Now he works with the Dept. of Telecommunications, Brno University of Technology.